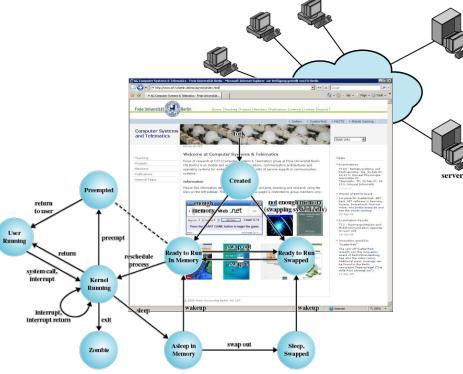


TI III: Operating Systems & Computer Networks Example

Prof. Dr.-Ing. Jochen Schiller Computer Systems & Telematics Freie Universität Berlin, Germany





Content

- 8. Networked Computer & Internet
- 9. Host-to-Network
- 10. Internetworking
- 11. Transport Layer
- 12. Applications
- 13. Network Security
- 14. Example

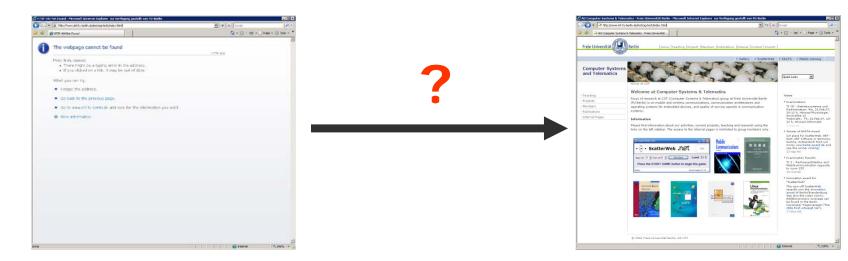


A Comprehensive Example

What happens if one presses a key on the computer?



What if that key causes an web page to be displayed?





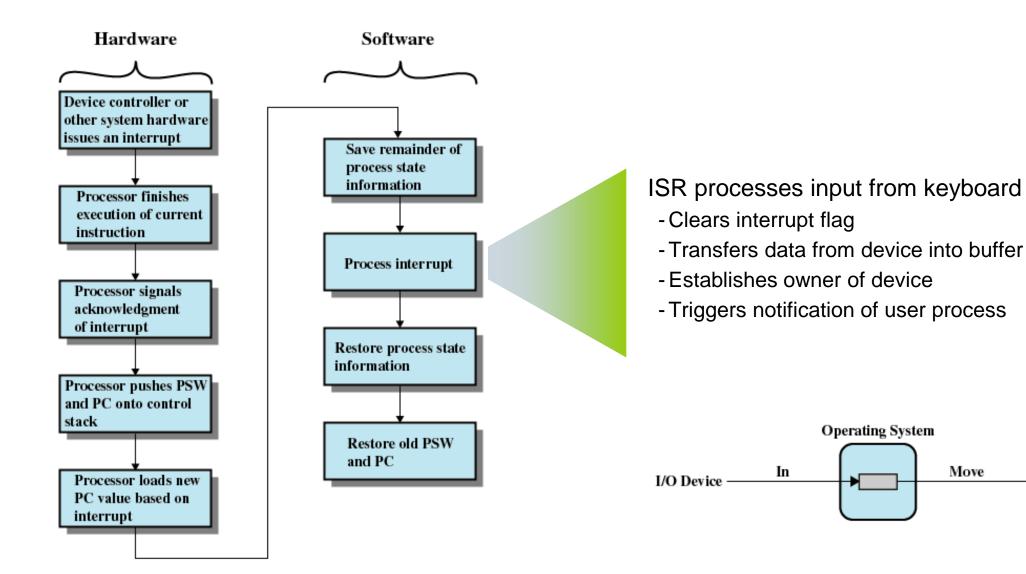
Keyboard Interrupt ΤĐΜ Control Keyboard controller raises interrupt flag Stack CPU interrupts execution of current process and starts Interrupt Т N+1 + Service Routine (ISR) Program Counter - Unconditional jump Y Start General Interrupt Registers Service Fetch Stage Execute Stage Interrupt Stage Routine Y+L Return $T \blacktriangleleft$ Stack Pointer Interrupts Disabled Processor Check for Fetch next Execute interrupt; START ΤĐΜ instruction initiate interrupt instruction Interrupts handler N N+1 Enabled User's Program HALT Main Memory



User Process

Move

Keyboard Interrupt Handling

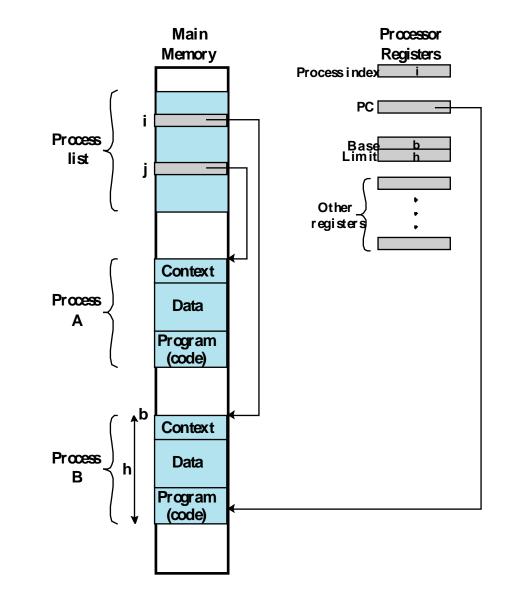




In the Meantime...

- Web browser is one of many processes running locally
- Other processes include
- -Other user processes (possibly of different users)
- System processes implementing system services
- -Kernel processes

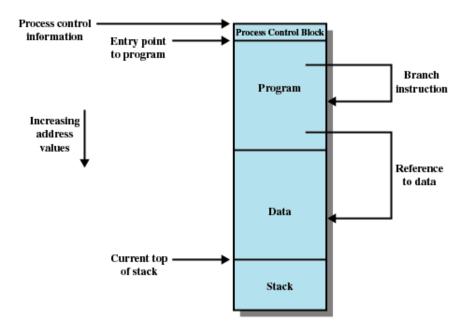
Ap	plications Processes	Performance Netw	orking	Users		
	Image Name	User Name	CPU	Mem Usage	^	
	System Idle Process	SYSTEM	94	20 K		
	taskmgr.exe	ed.bott	05	3,376 K		
	explorer.exe	ed.bott	01	10,696 K		
	notepad.exe	ed.bott	00 00	1,816 K		
	VMwareService.exe	SYSTEM		1,032 K		
	spoolsv.exe	SYSTEM	00	3,596 K		
	svchost.exe	LOCAL SERVICE	00	3,540 K		
	svchost.exe	NETWORK SERVICE	00	2,584 K		
	svchost.exe	SYSTEM	00	14,508 K		
	SnagIt32.exe	ed.bott	00	2,544 K		
	svchost.exe	SYSTEM	00	3,116 K		
	wuauclt.exe	ed.bott	00	3,752 K		
	lsass.exe	SYSTEM	00	984 K		
	services.exe	SYSTEM	00	2,460 K		
	winlogon.exe	SYSTEM	00	2,508 K		
	csrss.exe	SYSTEM	00	2,768 K	_	
	smss.exe	SYSTEM	00	348 K		
	msmsgs.exe	ed.bott	00	1,924 K		
	VMwareTrav.exe	ed.hott	nn	1.880 K		
	Show processes fro	End Process				

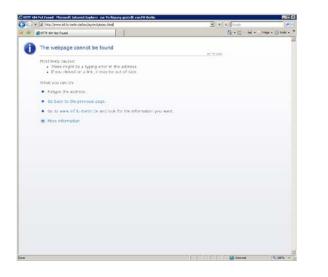


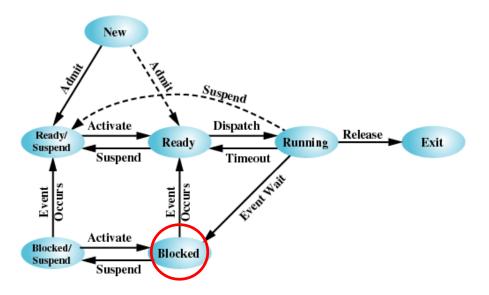


Web Browser Process in Detail

Web browser processes >Currently waiting for input -E.g. using select() >Process state blocked



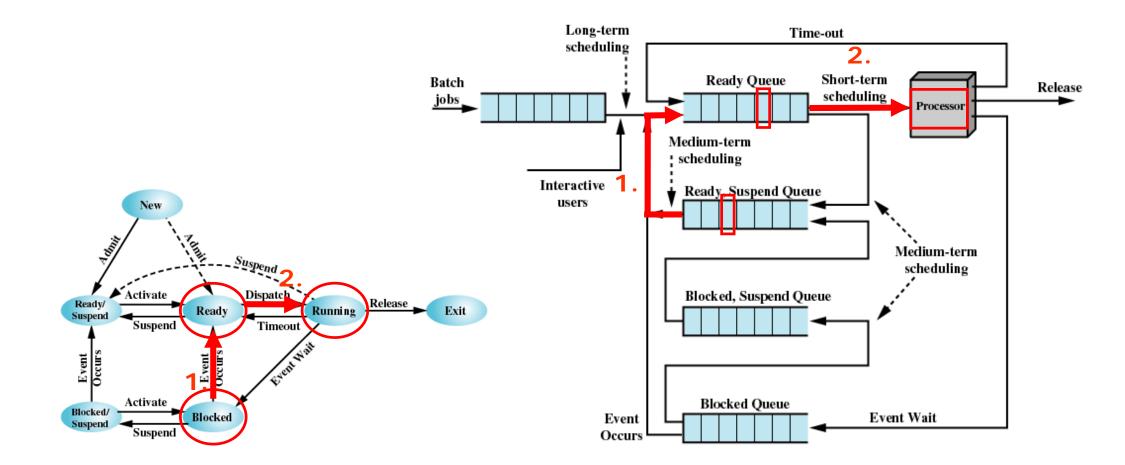






Reaction to External Event

- 1. ISR changes process state to ready
- 2. Scheduling algorithm eventually changes process state to running



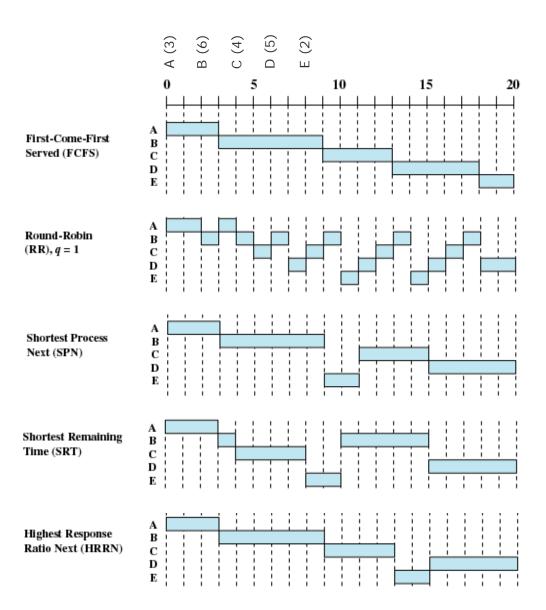


Process Scheduling

- Scheduling is handled by variety of scheduling algorithms
- -Non-preemptive / preemptive
- -Maximize throughput, responsiveness, etc...

Processes may have priorities

- Priority inversion due to lock on shared resources
- Priority inheritance

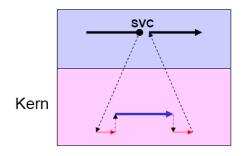




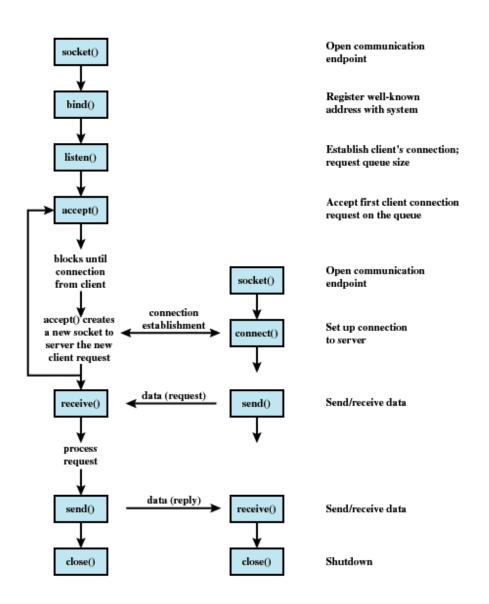
Web Browser Processes Event

Assume input requires web browser to display a web page with a given URL

- 1. String processing (user space)
- 2. Connect to server and retrieve necessary data (system calls)

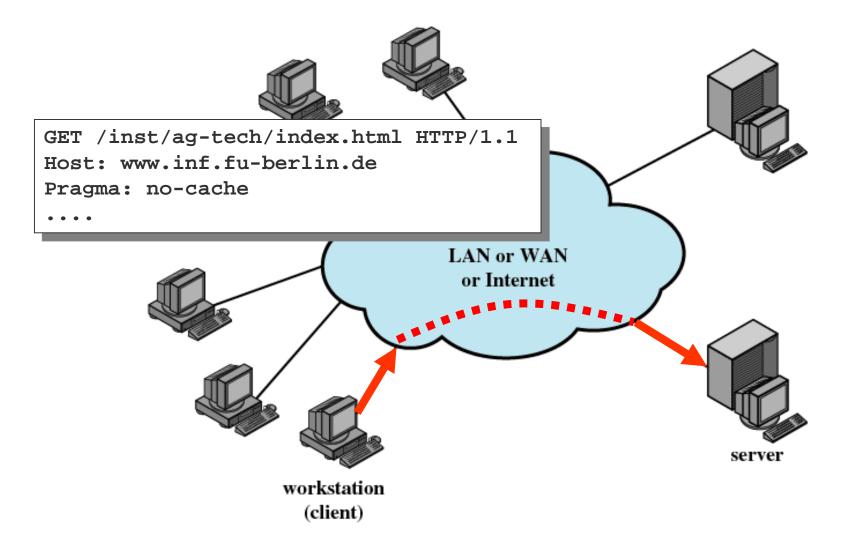


- 3. Render web page (user space)
- 4. Update user interface (system calls)



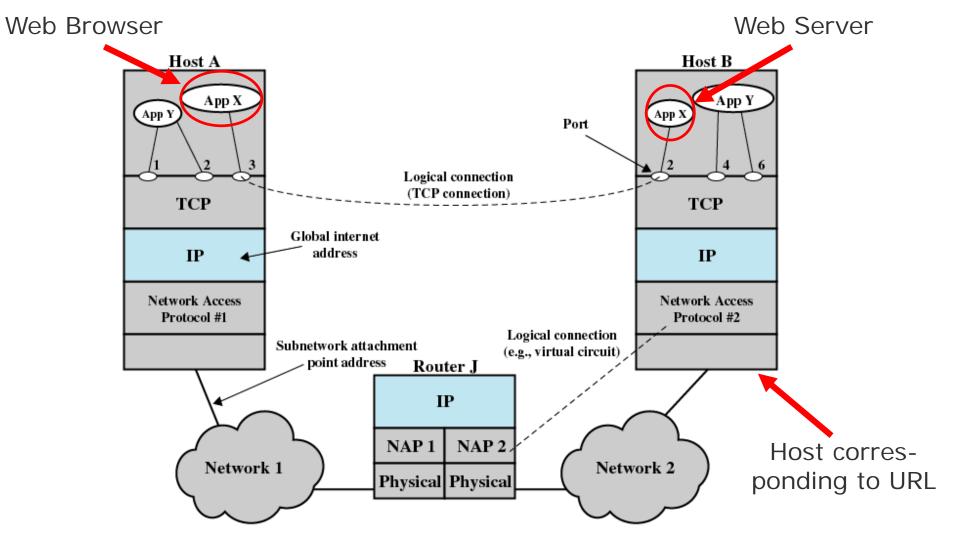


Client/Server Communication





Layered Protocol Stack





Interaction Between Network Layers

Layered protocol architecture

- Each layer uses only services of layer directly below

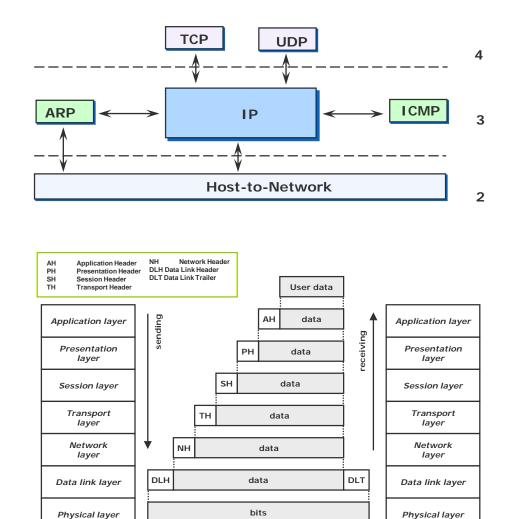
- Each layer provides services to layer directly above

➢Protocol independence

➤Modularity

Data encapsulation

- -Lower layers treat upper layer packets as simple data
- -Headers contain control information for each layer
- Repeated encapsulation causes overhead





Uniform Resource Locator (URL)

http://cst.mi.fu-berlin.de/index.html

http: Hypertext Transfer Protocol (HTTP)

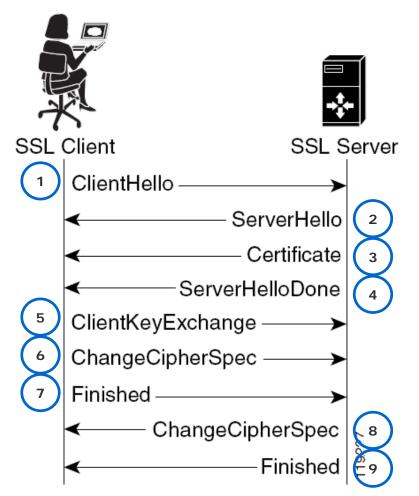
- Protocol for accessing web pages and related content
- Implies communication over port 80 (unless other port given in URL)
- cst.mi.fu-berlin.de: Host name
- -Resolved to IP address via Domain Name System (DNS)
- cst.mi.fu-berlin.de -> 160.45.117.167
- index.html: Local resource name
- Protocol specific parameter
- -Handled by web server



Security: HTTP over TLS/SSL

HTTPS authenticates server and establishes secure connection:

- 1) Propose SSL parameters, send random number
- 2) Agree to parameters, send random number
- 3) Send public key certificate
- 4) Conclude handshake negotiation
- 5) Send random number encrypted with server's public key
 - Client and server derive session key from all three random numbers
- 6) Activate negotiated parameters
- 7) Send encrypted hash over previous messages
 - Server decrypts and verifies message
- 8) Activate negotiated parameters
- 9) Send encrypted hash over previous messages
 - Client decrypts and verifies message
- Proceed to exchange regular HTTP data over secure channel



Source: Cisco Systems. Application Control Engine Module SSL Configuration Guide



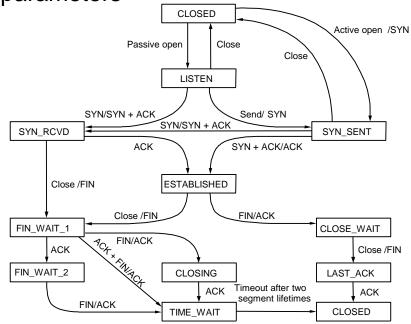
Connection Setup / Transport Layer

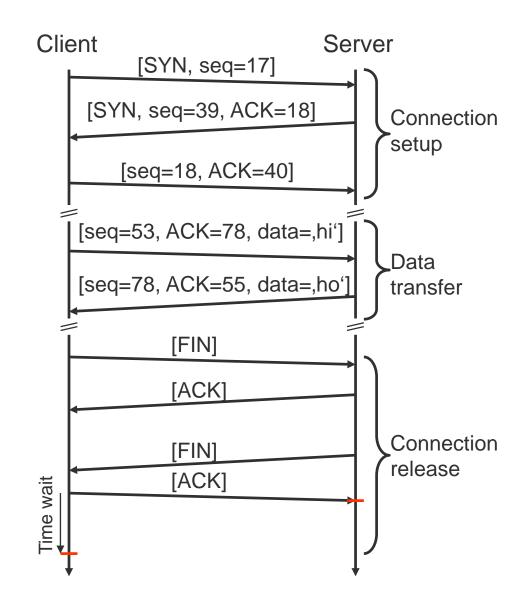
Reliable end-to-end connection between processes

Call to connect() initiates connection setup

➤TCP 3-way handshake

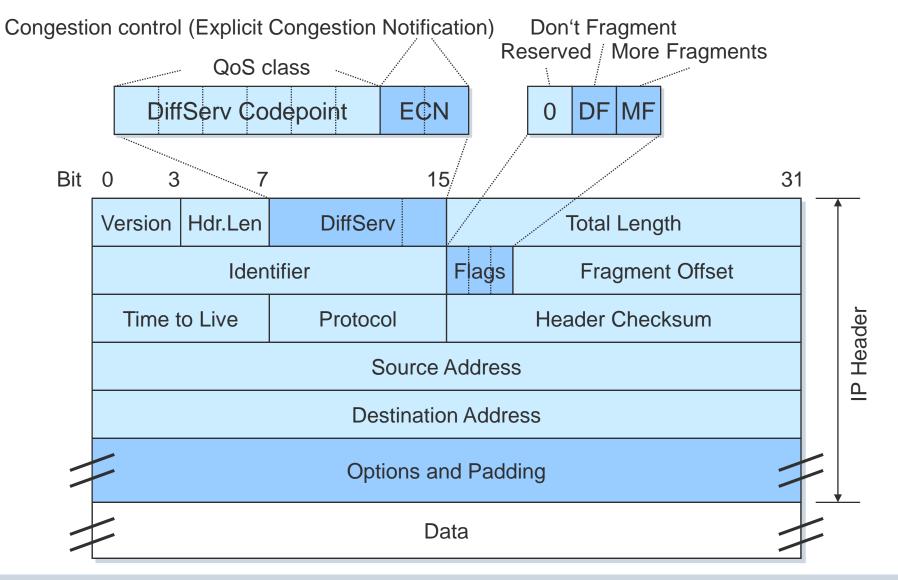
➤Connection parameters







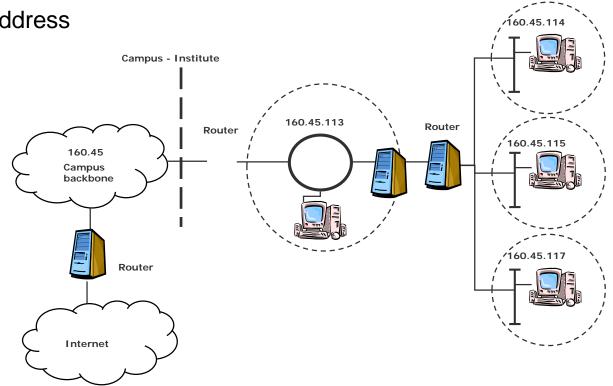
Structure of Network Layer IP-Packet





Network Layer Routing (Local Scope)

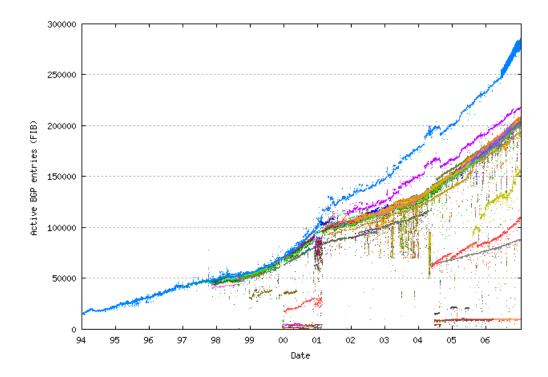
- Globally unique per host addressing
- Routers maintain tables of known networks
- -Optional route to default gateway
- Subnetting implements logical structure
- -Subnet mask builds hierarchy using host part of IP address
- -Limits broadcasts
- -More efficient routing
- Network topology may be part of security concept

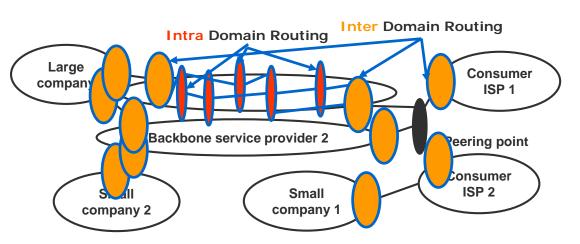




Network Layer Routing (Global Scope)

Internet organized into autonomous systems (AS) - Commonly, one AS per major organization - Peering points to exchange data between ASs Intra-domain routing: OSPF, link state algorithm Inter-domain routing: BGPv4, distance vector protocol - May involve non-technical routing choices



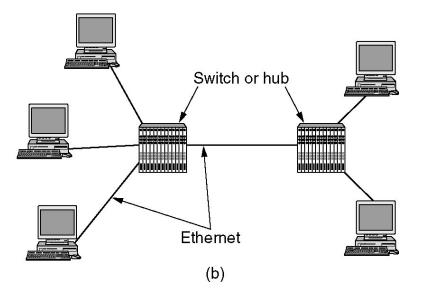


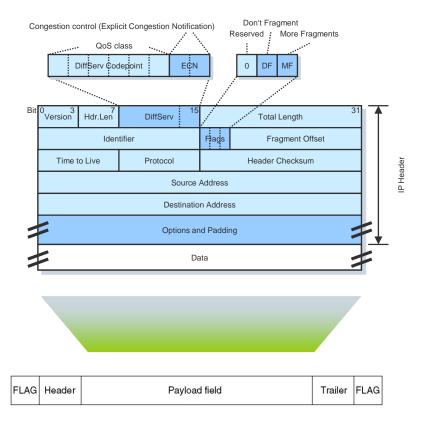


Data Link Layer Communication (Local Scope)

Transparent communication between two directly connected nodes

Services include: framing, error control, connection maintenance, acknowledgements, flow control







Error Detection: Cyclic Redundancy Check (CRC)

• Reception of a correct bit sequence:

```
11 0011 1001 ÷ 1 1001 = 10 0001 (mod 2)

<u>11 001</u>

00 0001 1001

<u>1 1001</u>

0 0000 = remainder
```

- No remainder, thus the received bits should be error free
- Reception of a erroneous bit sequence:

```
11 1111 1000 ÷ 1 1001 = 10 1001 (mod 2)

<u>11 001</u>

00 1101 1

<u>1100 1</u>

0001 0000

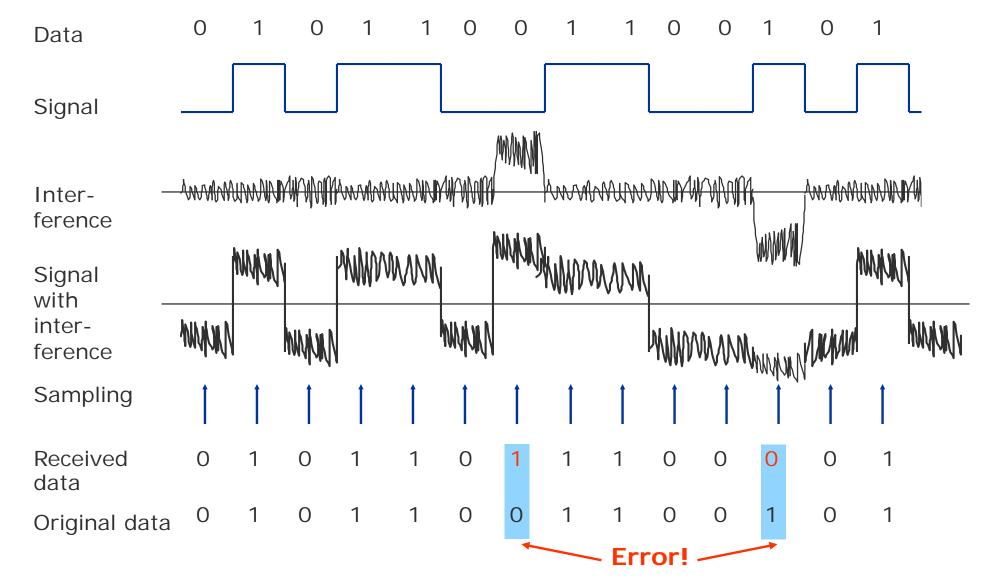
<u>1 1001</u>

0 1001 = remainder \neq 0
```

• There is a remainder unequal 0, thus there was definitely a transmission error



Errors During Transmission





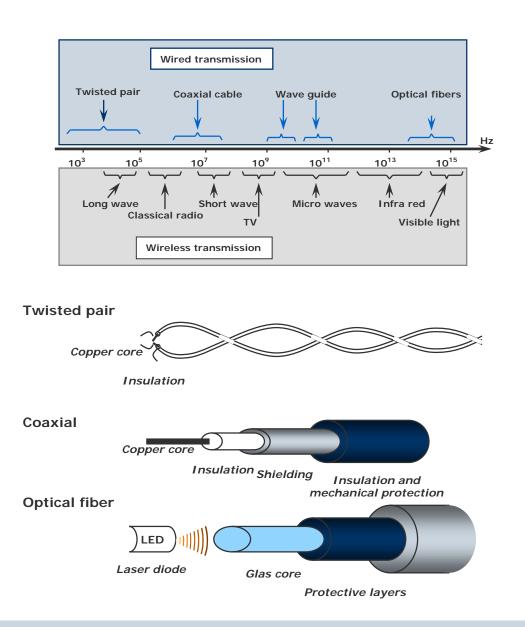
Physical Layer

Packet / sequence of bits turned into physical signal

Signal propagation depends on physical medium (limited bandwidth, attenuation, dispersion) and background noise

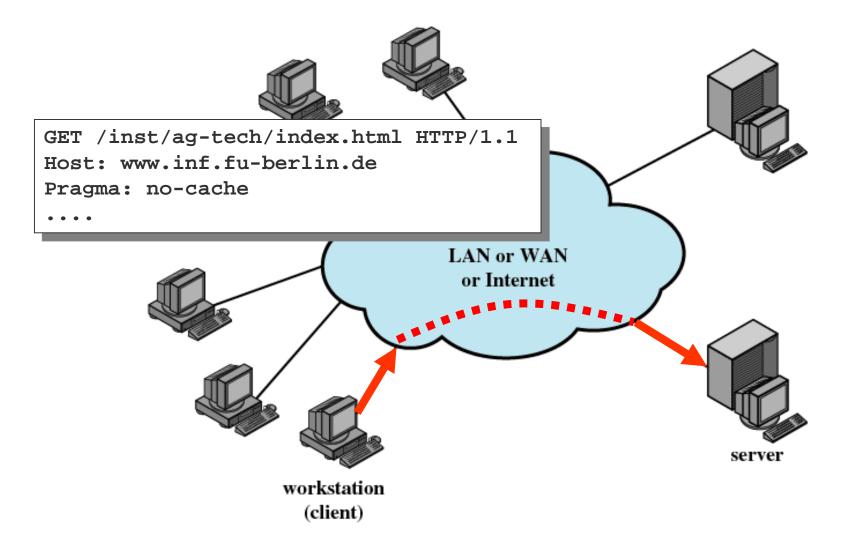
Mapping between bits and (multi-valued) symbols

Baseband transmission vs. modulation (broadband transmission)





Client/Server Communication





At the Server...

locally

	- 10:33:30 up 2 days, 1:04, 1 user, load average: 0.41, 0.26, 0. ks: 93 total, 1 running, 92 sleeping, 0 stopped, 0 zombie (s): 0.3% us, 0.0% sy, 0.0% ni, 99.7% id, 0.0% wa, 0.0% hi, 0.											
Cpu(s) Mem:										θ.8% ωα, θ.8% hi, θ free, 168112k buffe		
	2015992k											196k cached
	USER	PR		VIRT	RES			×CPU				Command
	wittenbu	16	0	1944	968	740		0.3				
	root	15	0	1584	520			0.0	0.1			init
	root	RT	0	0	0			0.0	0.0			migration/
	root	34		0	0			0.0	0.0			ksoftirqd∕
	root		-5	0	0			0.0	0.0			events/0
	root	13	-5		0			0.0	0.0			khelper
	root		-5	0	0			0.0	0.0			kthread
	root		-5		0			0.0	0.0			kblockd∕0
	root	10		0	0			0.0	0.0			khubd
	root		-5	0	0			0.0	0.0			kseriod
	root	20	0	0	0			0.0	0.0			pdflush
	root	15	0	0	0			0.0	0.0			pdflush
	root	15	0	0	0			0.0	0.0			kswapd0
107	root	20		0	0	0	S	0.0	0.0			aio/0
	root		-5	0	0			0.0	0.0			xfslogd∕0
109	root	20	-5	0	0	0	S	0.0	0.0	0:00	.00	xfsdatad/0
764	root	11	-5	0	0	0	S	0.0	0.0	0:00	.00	ata/0
781	root	11	-5	0	0	0	S	0.0	0.0	0:00	.00	kpsmoused

Web server is one of many processes running

 Upon receiving packet, network interface controller (NIC) will raise interrupt
 Kernel will handle the packet and notify the web server process

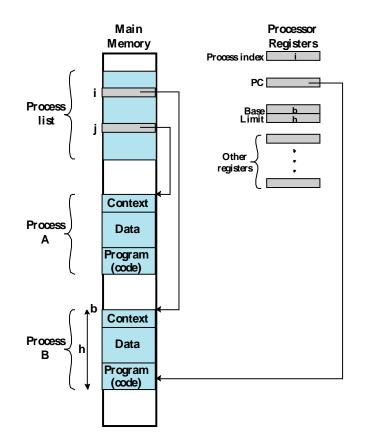
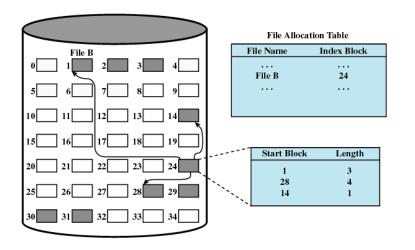


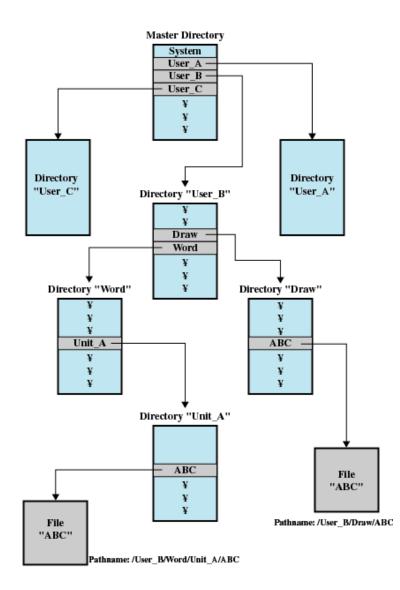
Figure 2.8 Typical Process Implementation



Processing of HTTP-GET Request

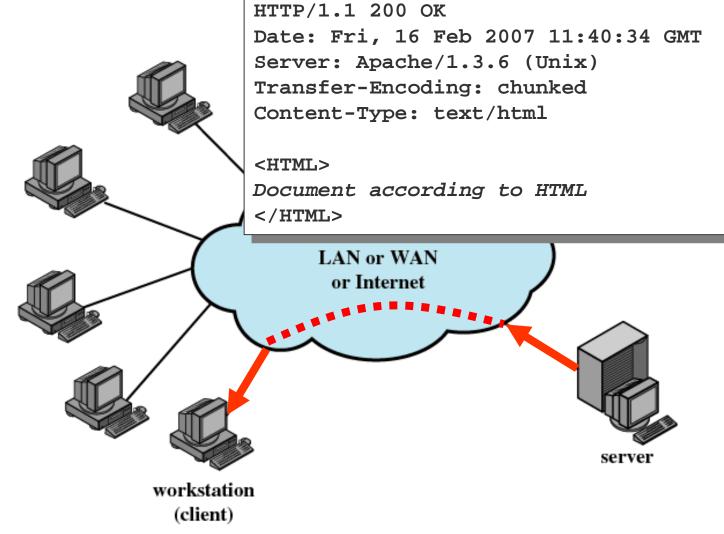
Web server retrieves file inst/agtech/index.html from local file system >System calls to access secondary storage >Kernel maps file name to data layout on disk Web server sends data to client







Server Replies to Client



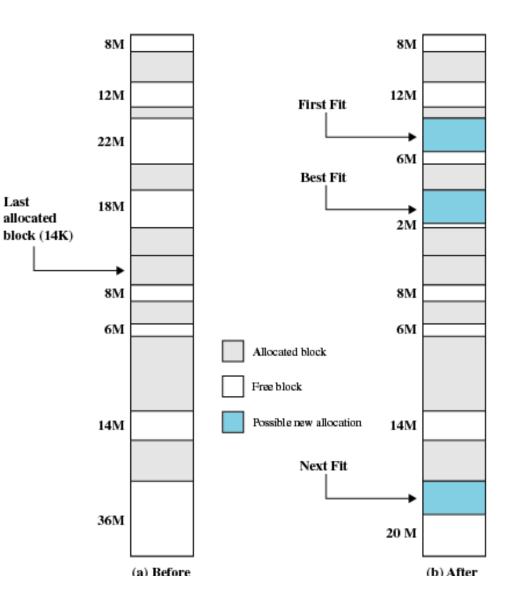


Client Data Processing

Client host receives packet Kernel hands data to web browser process Web browser renders page

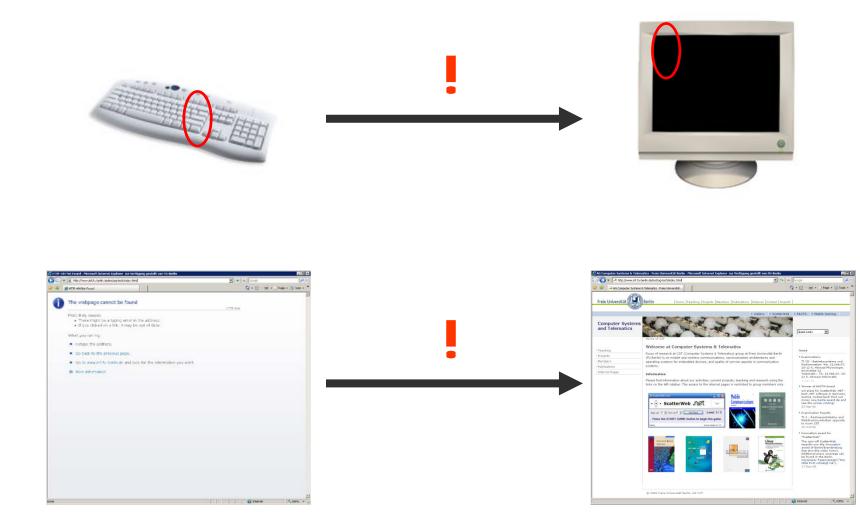
May have to allocate memory in the process
 Finally, browser updates user interface via system call







A Comprehensive Example



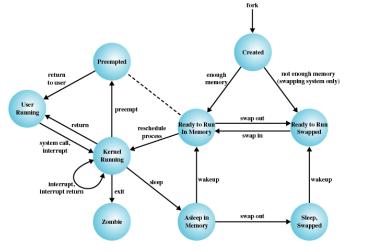


Content

- 1. Introduction and Motivation
- 2. Subsystems, Interrupts and System Calls
- 3. Processes
- 4. Memory
- 5. Scheduling
- 6. I/O and File System
- 7. Booting, Services, and Security

- 8. Networked Computer & Internet
- 9. Host-to-Network
- 10. Internetworking
- 11. Transport Layer
- 12. Applications
- 13. Network Security
- 14. Example







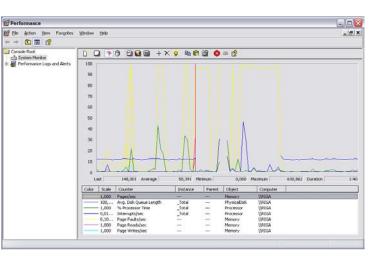


Figure 3.17 UNIX Process State Transition Diagram



